

## heritage - solution

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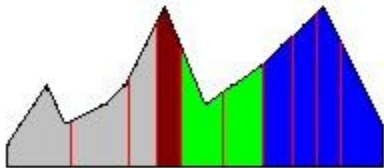
**First step.** If  $S$  is the sum of the  $n$  sons' age, we will parcel the polygon with  $S-1$  vertical fences in  $S$  equal areas. For example if the son has 4 sons of 1,2,3, respectively 4 years old, we have  $S=10$  so we parcel the polygon in 10 equal areas using 9 fences



**Second step.** Now we can observe that each son will take  $\lfloor S/i \rfloor$  consecutive areas. For example the sons could occupy the areas in the order [1 year, 2 years, 3 years and 4 years]. We will use the fences **1,3** and **6**.



Another example: the sons could occupy the areas in the order [3 years, 1 year, 2 years and 4 years]. We will use the fences **3,4** and **6**.



It doesn't matter the order in which the sons will occupy the areas from left to right, we will use  $n-1$  from the  $S-1$  fences we have already determined. Using all the possible orders we will have  $n! = 1*2*...*n$  cases. For each case we can determine the sum of the fences' length. We choose the smallest one.

### The algorithm's complexity



First step can be solved linearly (in  $O(m+S)$ ) if we pass from left to right through initial rectangular trapezoids and we use geometrical formulae in fixing the S-1 fences. We can also use binary search in this step.

Second step (backtracking) can have a  $O(n!)$  complexity enough to respect the time limit because for  $n = 8$  we have  $n! = 40320$ .